

• def: (successor) Let  $A$  be a set. The successor of  $A$  is the set  $A^+ = A \cup \{A\}$ .

• def: (Inductive set) A set  $\omega$  is a inductive (or successor) set iff  $\emptyset \in \omega$  & if  $n \in \omega$  then  $n^+ \in \omega$ .

• Axiom (in Zermelo-Fraenkel Set theory): There is an inductive set.

Problem 2: Let  $A$  be an inductive set.

(i) claim: if  $B$  is an inductive set then  $B \cap A \neq \emptyset$ .

PF  $A, B$  are inductive sets  $\therefore \emptyset \in A \wedge \emptyset \in B$  (by def.)

$\therefore \emptyset \in A \cap B$  by def. of intersection.

$\therefore A \cap B$  contains at least one element

$\therefore A \cap B$  is not empty  $\Rightarrow \emptyset \neq A \cap B$   $\square$

(ii) claim If  $\mathcal{I}$  is a family of inductive sets then

$\bigcap \mathcal{I}$  is an inductive set.

PF By definition  $\bigcap \mathcal{I} = \{x \mid x \in A \forall A \in \mathcal{I}\}$ .

Let  $A \in \mathcal{I}$ .  $A$  inductive  $\therefore \emptyset \in A$

$\therefore \forall A \in \mathcal{I}, \emptyset \in A$ .

$\therefore$  by definition  $\emptyset \in \bigcap \mathcal{I}$ .

Let  $n \in \bigcap \mathcal{I}$ . By definition  $\forall A \in \mathcal{I} n \in A$ .

Let  $A \in \mathcal{I}$ . Then  $n \in A$ .  $A$  inductive.

$\therefore n^+ \in A$ .  $\therefore \forall A \in \mathcal{I} n^+ \in A$

$\therefore$  by def.  $n^+ \in \bigcap \mathcal{I} \therefore \bigcap \mathcal{I}$  is an inductive set  $\square$

(2)

(iii) Define  $\tilde{\omega} = \bigcap \{B \subseteq A : B \text{ is inductive}\}$   
 $= \{x \mid x \in B \forall B \subseteq A : B \text{ inductive}\}$

claim: Let  $C$  be an inductive set. Then  $\tilde{\omega} \subseteq C$ .

PF By (ii)  $\tilde{\omega}$  is inductive. By (ii)  $A \cap C$  is inductive. (take your family  $\mathcal{I}$  to be  $\{A, C\}$ , then  $\bigcap \mathcal{I} = A \cap C$ .)

Let  $x \in A \cap C$ . Then  $x \in A \wedge x \in C$  by def of  $\cap$ .  
 $\therefore A \cap C \subseteq A \wedge A \cap C \subseteq C$  by def of  $\subseteq$ .  
Let  $x \in \tilde{\omega}$ . Then  $x \in B \forall B \subseteq A \wedge B$  inductive.  
 $A \cap C \subseteq A \wedge A \cap C$  inductive  $\therefore x \in A \cap C$ .  
Since  $A \cap C \subseteq C$ ,  $x \in C$ .  $\therefore \forall x \in \tilde{\omega}, x \in C$ .  
 $\therefore \tilde{\omega} \subseteq C$ .  $\square$

Problem #3:

Claim: IF  $P$  is linearly ordered  $\wedge f: P \rightarrow P$  is strictly increasing (i.e.  $p < q \Rightarrow f(p) < f(q)$ ) then  $f$  is 1-1.

PF let  $p \neq q \in P$ . Assume  $f(p) = f(q)$ .  
Since  $P$  is linearly ordered,  $p < q$  or  $q < p$  or  $p = q$ .

case 1: Assume  $p < q$ . Then  $f(p) < f(q)$   
Since  $f(p) = f(q)$ ,  $f(p) < f(p)$   
Contradiction.

Case 2: Assume  $q < p$ . Then  $f(q) < f(p)$

Since  $f(q) = f(p)$ ,  $f(q) < f(q)$

Contradiction.

$\therefore p = q$ .  $\therefore f$  is 1-1.  $\square$ .

Problem 4: Let  $\{P_i\}$  be the set of all ponds that have ever existed.

Let  $P_i(t)$  be the set of ducks on pond  $P_i$  at time  $t$ . Let

$R = \{ (x, y) \in D \mid \exists i, t \text{ s.t. } x \in P_i(t) \wedge y \in P_i(t) \}$ .

(i) reflexive: WTS  $(x, x) \in R \forall x \in D$ .

$\therefore$  WTS  $\exists i, t$  s.t.  $x \in P_i(t)$ . A priori no such  $i, t$  may exist. (if they did then  $R$  is reflexive)  
Not possible.

(ii) symmetric: Assume  $(x, y) \in R$ .  $\therefore \exists i, t$  s.t.  $x \in P_i(t) \wedge$

$y \in P_i(t) \Rightarrow \exists i, t$  s.t.  $y \in P_i(t) \wedge x \in P_i(t)$

$\therefore (y, x) \in R$ .

By definition  $R$  antisymmetric  $\Leftrightarrow$  (if  $(x, y) \in R \wedge (y, x) \in R$  then  $x = y$ ). Since  $R$  is symmetric,  $(x, y) \Leftrightarrow (y, x)$ .  $\therefore \forall i, \forall t$  if  $x, y \in P_i(t)$  then  $x = y$ . Who knows? Not possible.

(iii) transitive: Assume  $(x, y) \wedge (y, z) \in R$ .

then  $\exists i, t \wedge i', t'$  s.t.  $x, y \in P_i(t)$

$\wedge y, z \in P_{i'}(t')$ . But doesn't imply

$\exists i'', t''$  s.t.  $x, z \in P_{i''}(t'')$ .

(4)

Problem #5: Let  $A \subseteq \mathbb{Q}$  be an infinite set. Claim:

$\exists$  bijection  $f: \mathbb{Q} \rightarrow A$

PF  $\mathbb{Q}$  countable  $\therefore \exists$  bijection  $g: \mathbb{Q} \xrightarrow{\sim} \mathbb{N}$ .

Let  $g(A) = \{g(a) \mid a \in A\}$ .

Let  $g|_A: A \rightarrow g(A)$  be the function  $g|_A(a) = g(a)$ .

Then  $g|_A$  is a bijection ( $g|_A$  is 1-1;  $g|_A$  is clearly onto).

$\therefore g(A)$  is an infinite, hence unbounded subset of  $\mathbb{N} \therefore \exists$  a bijection  $h: \mathbb{N} \rightarrow g(A)$ .

Since  $g|_A$  is a bijection its inverse  $g|_A^{-1}: g(A) \rightarrow A$  is a bijection.

$\therefore$  We have the composition of bijections:

$$\mathbb{Q} \xrightarrow{g} \mathbb{N} \xrightarrow{h} g(A) \xrightarrow{g|_A^{-1}} A$$

Let  $f = g|_A^{-1} \circ h \circ g: \mathbb{Q} \rightarrow A$ .

The composition of bijections is a bijection.

$\square$